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10/642,360	-	08/15/2003	Adam G. Wolff	LZLO-01008US0	1630
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SAN FRANCISCO, CA 94105				ART UNIT	PAPER NUMBER
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Please find below and/or attached an Office communication concerning this application or proceeding.

	· · · · · · · · · · · · · · · · · · ·	Application No.	Applicant(s)
		10/642,360	WOLFF ET AL.
	Office Action Summary	Examiner	Art Unit
		Zheng Wei	2192
Period fo	The MAILING DATE of this communication apport	pears on the cover sheet with the c	orrespondence address
A SH WHIC - Exte after - If NC - Failu Any	ORTENED STATUTORY PERIOD FOR REPL CHEVER IS LONGER, FROM THE MAILING D insions of time may be available under the provisions of 37 CFR 1.1 SIX (6) MONTHS from the mailing date of this communication. O period for reply is specified above, the maximum statutory period ure to reply within the set or extended period for reply will, by statute reply received by the Office later than three months after the mailin led patent term adjustment. See 37 CFR 1.704(b).	ATE OF THIS COMMUNICATION 136(a). In no event, however, may a reply be tin will apply and will expire SIX (6) MONTHS from e, cause the application to become ABANDONE	N. nely filed the mailing date of this communication. D (35 U.S.C. § 133).
Status			
2a)□	Responsive to communication(s) filed on <u>15 A</u> This action is FINAL . 2b) This Since this application is in condition for alloward closed in accordance with the practice under the	s action is non-final. ince except for formal matters, pro	
Disposit	ion of Claims		
5)□ 6)⊠ 7)□	Claim(s) 1-37 is/are pending in the application 4a) Of the above claim(s) is/are withdra Claim(s) is/are allowed. Claim(s) 1-37 is/are rejected. Claim(s) is/are objected to. Claim(s) are subject to restriction and/or	wn from consideration.	
Applicat	ion Papers		•
10)⊠	The specification is objected to by the Examine The drawing(s) filed on <u>15 August 2003</u> is/are: Applicant may not request that any objection to the Replacement drawing sheet(s) including the correct The oath or declaration is objected to by the Examine	a)⊠ accepted or b)☐ objected drawing(s) be held in abeyance. Settion is required if the drawing(s) is ob	e 37 CFR 1.85(a). jected to. See 37 CFR 1.121(d).
Priority (under 35 U.S.C. § 119		
a)	Acknowledgment is made of a claim for foreign All b) Some * c) None of: 1. Certified copies of the priority document 2. Certified copies of the priority document 3. Copies of the certified copies of the priority application from the International Burea See the attached detailed Office action for a list	ts have been received. ts have been received in Applicationity documents have been receive u (PCT Rule 17.2(a)).	on No ed in this National Stage
Attachmen	it(s)		
2) D Notic 3) D Infor	ce of References Cited (PTO-892) ce of Draftsperson's Patent Drawing Review (PTO-948) mation Disclosure Statement(s) (PTO/SB/08) er No(s)/Mail Date 08/15/2003.	4) Interview Summary Paper No(s)/Mail Di 5) Notice of Informal F 6) Other:	ate

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DETAILED ACTION

1. This office action is in response to the application filed on 08/15/2003.

2. Claims 1-37 are pending and have been examined.

Priority

3. The priority date for this application is 08/15/2003. No continuing data and foreign applications are related to this application.

Information Disclosure Statement

4. The information disclosure statement filed 08/15/2003 has been placed in the application file and the information referred to therein has been considered.

Claim Rejections - 35 USC § 102

5. The following is a quotation of the appropriate paragraphs of 35 U.S.C. 102 that form the basis for the rejections under this section made in this Office action:

A person shall be entitled to a patent unless -

(b) the invention was patented or described in a printed publication in this or a foreign country or in public use or on sale in this country, more than one year prior to the date of application for patent in the United States.

6. Claims 1-5, 7, 13-15, 18, 19, 21, 26, 27, 29 and 35 are rejected under 35

U.S.C. 102(b) as being anticipated by <u>Steele</u> (Guy L. Steele, Common Lisp the language, 2nd edition)

Claim 1:

<u>Steele</u> discloses a method to provide for evaluating expressions, comprising:

- Receiving code for a program, said code includes one or more
 expressions and one or more markers that specify when said one or more
 expressions should be evaluated during execution of said program (see
 for example, section 5.3.3, Control of Time of Evaluation, [Special Form],
 p.11, line 11, "situation" and related text); and
- Automatically providing additional functionality to said code for said program, said additional functionality evaluates said one or more expressions during execution of said program at one or more times specified by said one or more markers (see for example, section 5.3.3, Control of Time of Evaluation, p.11, line 14, "interpreter", "compiler"),

Claim 2:

<u>Steele</u> further discloses a method according to claim 1, wherein: said one or more markers specify when said one or more expressions should be evaluated during execution of said program independent from a context of where said expressions are used (see for example, section 5.3.3, Control of Time of

Evaluation, p.11, line 8, "...executed only at compile time, only at load time or when interpreted but not compiled.")

Claim 3:

Steele also discloses a method according to claim 1, wherein: said markers can indicate that a particular expression should be evaluated immediately, once or always (see for example, section 5.3.3, Control of Time of Evaluation, p.11, lines 12-15, "Each *situation* must be a symbol, either *compile*, *load or eval. eval* specifies that the interpreter should process the body. *compile* specifies that compiler should evaluate the body at compile time in the compilation context").

Claim 4:

Steele further discloses a method according to claim 3, wherein: failure of a marker to indicate that a particular expression should be evaluated immediately or once defaults to an indication indicate that said particular expression should be evaluated always (see for example, section 5.3.3, Control of Time of Evaluation, p.11, lines 32-34, "If the form is not an eval-when form...perform normal compiler processing of the form").

Claim 5:

Steele further discloses a method according to claim 1, wherein:

- Said one or more expressions are constraints for variables (see for example, section 5.3.3, Control of Time of Evaluation, p.12, lines 1-4, example code, "lambda (stream char)", "declare (ignore char)"; and
- Said step of automatically providing additional functionality to said code includes adding codes that creates an object for each constraint, adds functions to said object that sets said variables, and adds functions that sets dependencies for said expressions (see for example, section 5.3.3, Control of Time of Evaluation, p.12, lines 1-5, example code and related text explanation, "This causes the call to set-macro-character to be executed in the compiler's execution environment, thereby modifying its reader syntax table.").

Claim 7:

<u>Steele</u> also discloses a method according to claim 1, wherein: said step of automatically providing additional functionality to said code includes compiling said code (see for example, section 5.3.3, Control of Time of Evaluation, p.11, lines 14-15, "*compile* specifies that compiler should evaluate the body at compile time in the compilation context").

Claim 13:

Steele discloses a method to provide for evaluating expressions, comprising:

- Receiving code for a program, said code includes one or more
 expressions and one or more markers that specify when said one or more
 expressions should be evaluated during execution of said program; (see
 for example, section 5.3.3, Control of Time of Evaluation, [Special Form],
 p.11, line 11, "situation" and related text) and
- Evaluating said one or more expressions during execution of said program
 at times specified by said one or more markers (see for example, section
 5.3.3, Control of Time of Evaluation, p.11, line 14, "interpreter",
 "compiler"),

Claim 14:

Steele further discloses a method according to claim 13, wherein: said one or more markers specify when said one or more expressions should be evaluated during execution of said program independent from a context of where said expressions are used (see for example, section 5.3.3, Control of Time of Evaluation, p.11, line 8, "... executed only at compile time, only at load time or when interpreted but not compiled.").

Claim 15:

<u>Steele</u> further discloses method according to claim 13, wherein: said markers can indicate that a particular expression should be evaluated immediately, once or always (see for example, section 5.3.3, Control of Time of Evaluation, p.11, lines

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12-15, "Each *situation* must be a symbol, either *compile*, *load or eval. eval* specifies that the interpreter should process the body. *compile* specifies that compiler should evaluate the body at compile time in the compilation context").

Claim 18:

Steele discloses a method to provide for evaluating expressions, comprising:

- Accessing code that includes an expression defining a first variable, said
 expression is dependent on a changeable item (see for example of
 implementation a *defun* form, p.14, line 14, a first variable "x" and
 expression); and
- Compiling said code, said step of compiling said code adds additional
 functionality to said code, said additional functionality evaluates said
 expression when said item changes and updates said first variable (see
 for example code of implementation to expand a defun form into evalwhen, p.14, line 16-25 and related text explanation).

Claim 19:

Steele further discloses a method according to claim 18, wherein:

 Said expression is part of a constraint for said first variable (see for example code of implementation to expand a defun form into eval-when, p.14, line 17, "compiler::notice-function 'bar ' (x)") Application/Control Number: 10/642,360 Page 8

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Said step of compiling includes creating an object for said constraint,
adding a first function to said object that sets said first variable,
determining dependency of said expression and adding a second function
for said dependency (see for example code of implementation to expand a
defun form into eval-when, p.14, line 16-25, "lambda (x)", "lambda () (+ x
3)" and related text explanation).

Claim 21:

Steele further discloses a method according to claim 18, wherein: said code includes a marker for said expression, said marker specifies when said expression should be evaluated during execution of said code (see for example code of implementation to expand a defun form into eval-when, p.14, lines 16-25, marker ":compile-toplevel" and related text explanation).

Claim 26:

<u>Steele</u> discloses a method to provide for evaluating expressions, comprising:

- Receiving code that includes an expression defining a first variable, said
 expression is dependent on a changeable item (see for example of
 implementation a defun form, p.14, line 14, a first variable "x" and
 expression); and
- Automatically providing additional functionality to said code, said
 additional functionality evaluates said expression when said item changes

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and updates said first variable (see for example code of implementation to expand a defun form into eval-when, p.14, line 16-25 and related text explanation).

Claim 27:

Steele further discloses a method according to claim 26, wherein:

- Said expression is part of a constraint for said first variable (see for example code of implementation to expand a defun form into eval-when, p.14, line 17, "compiler::notice-function 'bar ' (x)"); and
- Said step of automatically providing includes creating an object for said constraint, adding a first function to said object that sets said first variable, determining dependency of said expression and adding a second function for said dependency to said object (see for example code of implementation to expand a defun form into eval-when, p.14, line 16-25, "lambda (x)", "lambda () (+ x 3)" and related text explanation).

Claim 29:

Steele discloses a method according to claim 26, wherein: said code includes a marker for said expression, said marker specifies when said expression should be evaluated during execution of said code (see for example code of implementation to expand a defun form into eval-when, p.14, line 16-25, marker ":compile-toplevel" and related text explanation).

Claim 35:

<u>Steele</u> further discloses one or more processor readable storage devices according to claim 31, wherein: said preexisting additional functionality prevents circular evaluation (see for example, p.13, line 32, "It is never the case that the execution of a single *eval-when* expression will execute the body code more than once").

Claim Rejections - 35 USC § 103

- 7. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:
 - (a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negatived by the manner in which the invention was made.
- 8. Claims 6, 8, 16, 30 and 36 are rejected under 35 U.S.C. 103(a) as being unpatentable over <u>Steele</u> (Guy L. Steele, Common Lisp the language, 2nd edition) in view of <u>Rodriguez</u> (Rodriguez et al, Active Lisp Page)

Claim 6:

<u>Steele</u> discloses a method as in claim 1 above, but does not disclose said code for said program is XML code. However, <u>Rodriguez</u> in the same analogous art of using Lisp discloses an Active Lisp Pages which uses LISP in XML as a server scripting environment of creating dynamic pages and interactive applications

over the Web (see for example, p.3, right col., lines 34-45, "XML Syntax"). Therefore, it would have been obvious to one having ordinary skills in the art at the time the invention was made to use <u>Steele's LISP language</u> in ALP and further embed with well-formed XML as disclosed by <u>Rodriguez</u>. One would have been motivated to do so to take advantage of well-formed XML and ALP server page engine solution as once suggested by <u>Rodriguez</u>. (see for example, p.1, left col., line 31 – right col., line 53, "Motivation" of ALP and also see p.3, right col., lines 34-45, "XML Syntax", "The advantage of having well-formed XML is the use that one can make of many available tools for such format, such as validators, schemas, and XSLT stylesheets.")

Claim 8:

<u>Steele</u> discloses a method as in claim 1 above, but does not disclose the method for implementation of client-server service. However, <u>Rodriguez</u> in the same analogous art of method for using LISP in ALP discloses the method comprising:

- Receiving a request for content via a network (see for example, p.1, Abstract, lines 14-15, "a client requests an ALP file/page from the server.");
- Transmitting said code with said additional functionality to a client via said network (see for example, p.1, Abstract, lines 23-24, "The server transmits standard HTML to the browser."); and

 Executing said code with said additional functionality at said client (see for example, p.1, Abstract, lines 25-26, "... only the result if the script is returned to the browser.")

Therefore, it would have been obvious to one having ordinary skills in the art at the time the invention was made to use <u>Steele</u>'s LISP language in ALP to provide Web service. One would have been motivated to do so to provide secure service with simple, efficient advantages of ALP as once suggested by <u>Rodriguez</u>. (see for example, p.1, left col., line 31 – right col., line 53, "Motivation" of ALP)

Claim 16:

Steele discloses a method as in claim 13 above, but does not disclose said code for said program is XML source code. However, Rodriguez discloses an Active Lisp Pages which uses LISP in XML as a server side scripting environment of creating dynamic pages and interactive applications over the Web (see for example, p.3, right col., lines 34-45, "XML Syntax"). Therefore, it would have been obvious to one having ordinary skills in the art at the time the invention was made to use Steele's LISP language in ALP and further embed with well-formed XML as disclosed by Rodriguez. One would have been motivated to do so to take advantage of well-formed XML and ALP server page engine solution as once suggested by Rodriguez. (see for example, p.1, left col., line 31 – right col., line 53, "Motivation" of ALP and also see p.3, right col., lines 34-45, "XML Syntax", "The advantage of having well-formed XML is the use that one can

make of many available tools for such format, such as validators, schemas, and XSLT stylesheets.").

Claim 30:

Steele discloses a method as claim 26 above, but does not disclose the method for implementation of client-server service. However, Rodriguez in the same analogous art of method for using LISP in ALP discloses the method comprising:

- Requesting said code by an Internet client (see for example, p.1, Abstract, lines 14-15, "a client requests an ALP file/page from the server.");
- Transmitting said code with said additional functionality to said Internet client after said step of automatically providing (see for example, p.1, Abstract, lines 23-24, "The server transmits standard HTML to the browser."); and
- Executing said code with said additional functionality using said Internet client (see for example, p.1, Abstract, lines 25-26, "...only the result if the script is returned to the browser.").

Therefore, it would have been obvious to one having ordinary skills in the art at the time the invention was made to use <u>Steele</u>'s LISP language in ALP to provide Web service as disclosed by <u>Rodriguez</u>. One would have been motivated to do so to provide secure service with simple, efficient advantages of ALP as once

suggested by <u>Rodriguez</u>. (see for example, p.1, left col., line 31 – right col., line 53, "Motivation" of ALP).

Claim 36:

<u>Steele</u> discloses a method in LISP to evaluate expression comprising the steps of:

- accessing code that includes an expression defining a first variable, said expression is dependent on a changeable item (see for example of implementation a *defun* form, p.14, line 14, a first variable "x" and expression); and
- automatically providing preexisting additional functionality to said code, said preexisting additional functionality evaluates said expression when said item changes and updates said first variable (see for example code of implementation to expand a defun form into eval-when, p.14, line 16-25 and related text explanation).

But does not explicitly disclose said apparatus that provides these service.

However, Rodriguez in the same analogous art of method for using LISP in ALP discloses one or more processors in communication with said processor readable storage device, said one or more processors perform a method as Steele disclosed. (see for example, p.1, lines 11-13, "Web server", and "ALP processor"). Therefore, it would have been obvious to one having ordinary skills in the art at the time the invention was made to use Rodriguez's ALP server with

processor and storage device embodied to perform <u>Steele</u>'s method. One would have been motivated to do that, because it is efficient, simple and reusable. (see for example, p.1, left col., line 31 – right col., line 53, section Motivation, "Reusable components", "Efficiency", "Simplicity and fast development")

9. Claim 9-12, 22, 23, 25, 31, 32 and 34 are rejected under 35 U.S.C. 103(a) as being unpatentable over <u>Steele</u> (Guy L. Steele, Common Lisp the language, 2nd edition)

Claims 9-12:

Claims 9-12 are processor readable storage devices having processor readable code embodied on them, which are the product version of the claimed methods discussed as in claims 1-3 and 7 above. It is well known in the computer art to practice and store the computer readable code in such computer readable storage devices. Therefore, these claims are also obvious over <u>Steele</u>.

Claims 22, 23 and 25:

Claims 22, 23 and 25 are processor readable storage devices having processor readable code embodied on them, which are the product version of the claimed methods discussed as in claims 18, 19 and 21 above. It is well known in the computer art to practice and store the computer readable code in such computer readable storage devices. Therefore, these claims are also obvious over <u>Steele</u>.

Claims 31, 32 and 34:

Claims 31, 32 and 34 are processor readable storage devices having processor readable code embodied on them, which are the product version of the claimed methods discussed as in claims 26, 27 and 29 above. It is well known in the computer art to practice and store the computer readable code in such computer readable storage devices. Therefore, these claims are also obvious over <u>Steele</u>.

Claim 17 is rejected under 35 U.S.C. 103(a) as being unpatentable over <u>Steele</u>
 (Guy L. Steele, Common Lisp the language, 2nd edition) in view of <u>Hickey</u> (Hickey et al., LISP – a Language for Internet Script and Programming).

Claim 17:

Steele discloses a method as in claim 13 above, but does not disclose said code for said program is object code. However, Hickey discloses a method to implement LISP Applets which is object code and embed into web pages (see for example, p.7, section 3, LISP Applets, line 32 "Java byte code class file"). Therefore, it would have been obvious to one having ordinary skills in the art at the time the invention was made to use Steele's LISP language including special form eval-when to create LISP Applet for web pages. One would have been motivated to use LISP creating LISP applet object code which Hickey discloses to provide more efficient applets and greatly decreased download times and also

be able to make use of the latest Java compilation technology. (see for example, p.7, lines 37-39, "potential to provide more efficient applets and greatly decreased download times. Another advantage of this approach is that by compiling to Java, we are able to make use of the latest Java compilation technology.").

11. Claims 20, 24, 28 and 33 are rejected under 35 U.S.C. 103(a) as being unpatentable over <u>Steele</u> (Guy L. Steele, Common Lisp the language, 2nd edition) in view of <u>Haible</u> (Haible et al., Implementation Notes for GNU CLISP)
Claim 20:

Steele discloses a method as in claim 19 above, wherein: said additional functionality includes code that adds said first function to an object for said first variable, but does not explicitly disclose said code that provides a pointer to said first function to an object for said changeable item to be called by said object for said changeable item when said changeable item changes. However, Haible in the same analogous art of LISP implementation discloses that using "Weak Pointers" and "Foreign Pointers" to implement LISP in CLISP project (see for example, p.5, section: extensions 1.6, Weak Pointers, line 22, "A weak pointer is an object holding a reference to a given object"). Therefore, it would have been obvious to one having ordinary skills in the art at the time the invention was made to use Haible's weak pointer to implement Steele's example. One would have

been motivated to do so to quick access reference and without keeping the latter from being garbage-collected. (see for example, p.5, section: extensions 1.6, Weak Pointers, line 22, "without keeping the latter from being garbage-collected.")

Claim 24:

Claim 24 discloses processor readable code embodied on processor readable storage devices, which can be executed by processor to perform the method as in claims 20 above for evaluating expressions. Therefore, this claim is also unpatentable by <u>Steele</u> and <u>Haible</u>, as it is well known in the computer art to store such methods in a processor readable device.

Claim 28:

Steele discloses a method according to claim 27, wherein: said additional functionality includes code that adds said first function to an object for said first variable, but does not explicitly disclose said code that provides a pointer to said first function to an object for said changeable item to be called by said object for said changeable item when said changeable item changes. However, Haible in the same analogous art of LISP implementation discloses that using "Weak Pointers" and "Foreign Pointers" to implement LISP in CLISP project (see for example, p.5, section: extensions 1.6, Weak Pointers, line 22, "A weak pointer is an object holding a reference to a given object"). Therefore, it would have been

obvious to one having ordinary skills in the art at the time the invention was made to use <u>Haible</u>'s weak pointer to implement <u>Steele</u>'s example. One would have been motivated to do so to quick access reference and without keeping the latter from being garbage-collected. (see for example, p.5, section: extensions 1.6, Weak Pointers, line 22, "without keeping the latter from being garbage-collected.")

Claim 33:

Claim 33 discloses processor readable code embodied on processor readable storage devices, which can be executed by processor to perform the method as in claims 28 above for evaluating expressions. Therefore, this claim is also unpatentable by Steele and Haible, as it is well known in the computer art to store such methods in a processor readable device.

12. Claim 37 is rejected under 35 U.S.C. 103(a) as being unpatentable over <u>Steele</u>

(Guy L. Steele, Common Lisp the language, 2nd edition) in view of <u>Rodriguez</u>

(Rodriguez et al, Active Lisp Page) and further in view of <u>Haible</u> (Haible et al.,

Implementation Notes for GNU CLISP)

Claim 37:

<u>Steele</u> and <u>Rodriguez</u> disclose an apparatus according to claim 36 and <u>Steele</u> further discloses, wherein:

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Said expression is part of a constraint for said first variable (see for example code of implementation to expand a defun form into eval-when, p.14, line 17, "compiler::notice-function 'bar ' (x)");

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 Said step of automatically providing includes creating an object for said constraint, adding a first function to said object that sets said first variable, determining dependency of said expression and adding a second function for said dependency to said object (see for example code of implementation to expand a defun form into eval-when, p.14, line 16-25,

"lambda (x)", "lambda () (+ x 3)" and related text explanation); and
But both of them do not disclose said additional functionality includes code that
adds said first function to an object for said first variable and code that provides a
pointer to said first function to an object for said changeable item to be called by
said object for said changeable item when said changeable item changes.
However, <u>Haible</u> in the same analogous art of LISP implementation discloses
that using "Weak Pointers" and "Foreign Pointers" to implement LISP in CLISP
project (see for example, p.5, section: extensions 1.6, Weak Pointers, line 22, "A
weak pointer is an object holding a reference to a given object"). Therefore, it
would have been obvious to one having ordinary skills in the art at the time the
invention was made to use <u>Haible</u>'s weak pointer to implement <u>Steele</u>'s example
and use <u>Rodriguez</u>'s ALP server. One would have been motivated to do so to
quick access reference and without keeping the latter from being garbage-

collected. (see for example, p.5, section: extensions 1.6, Weak Pointers, line 22, "without keeping the latter from being garbage-collected.")

Conclusion

- 13. The prior art made of record and not relied upon is considered pertinent to applicant's disclosure.
 - Denis Attal (US 5860010) discloses the user of language with similar representation for programs and data.

Any inquiry concerning this communication or earlier communications from the examiner should be directed to Zheng Wei whose telephone number is (571) 270-1059. The examiner can normally be reached on Monday-Thursday 8:00-15:00.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Tuan Q. Dam can be reached on (571) 272-3695. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

Any inquiry of a general nature of relating to the status of this application or proceeding should be directed to the TC 2100 Group receptionist whose telephone number is 571- 272-1000.

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Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see http://pair-direct.uspto.gov. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free). If you would like assistance from a USPTO Customer Service Representative or access to the automated information system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.

ZW

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